

Minimal affinizations and their graded limits

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Representation theory and Related Topics @ Irako View Hotel

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Introduction

Jacobi-Trudi formula

For a partition $\lambda = (\lambda_1 \geq \cdots \geq \lambda_n)$,

$$s_\lambda(x) = \det (h_{\lambda_i - i + j}(x))_{1 \leq i, j \leq n}.$$

$s_\lambda(x)$: Schur polynomial, $h_k(x)$: complete symm. polynomial.

Translation in the \mathfrak{sl}_{n+1} -modules

$\lambda \in P^+$: dom. int. wt $\rightsquigarrow \lambda = (\lambda_1 \geq \cdots \geq \lambda_n)$ by $\lambda_i = \sum_{k \geq i} \langle h_k, \lambda \rangle$
 $\text{ch } V(\lambda) = s_\lambda(x)$, $\text{ch } V(k\varpi_1) = h_k(x)$ ($V(\lambda)$: simple \mathfrak{sl}_{n+1} -mod.)

Theorem

$$\text{ch } V(\lambda) = \det \left(\text{ch } V((\lambda_i - i + j)\varpi_1) \right)_{1 \leq i, j \leq n}$$

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So $\text{ch } V(\lambda) = \det \left(\text{ch } V((\lambda_i - i + j)\varpi_1) \right)_{1 \leq i, j \leq n}$ holds in type A.

Q. Does this formula hold in other types? **No!**

$$\text{ch } V(\lambda) \neq \det \left(\text{ch } V((\lambda_i - i + j)\varpi_1) \right)_{1 \leq i, j \leq n},$$

when $\mathfrak{g} \neq \mathfrak{sl}_{n+1}$ (though there may be several generalizations.)

Q. When $\mathfrak{g} \neq \mathfrak{sl}_{n+1}$, does $\det \left(\text{ch } V((\lambda_i - i + j)\varpi_1) \right)_{1 \leq i, j \leq n}$

have some representation theoretic meaning? **Yes!**

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In type BD , we have

$$\text{ch } L_q(\lambda) = \det \left(\text{ch } V((\lambda_i - i + j)\varpi_1) \right)_{1 \leq i, j \leq n},$$

where $L_q(\lambda)$ denotes a **minimal affinization** (a special class of f.d. simple $U_q(\mathcal{L}\mathfrak{g})$ -modules explained later).

In type C , a similar formula holds:

$$\text{ch } L_q(\lambda) = \det \left(\sum_{0 \leq 2k \leq \lambda_i - i + j} \text{ch } V((\lambda_i - i + j - 2k)\varpi_1) \right)_{1 \leq i, j \leq n}.$$

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2. Main Theorem (JT formula for $\text{ch } L_q(\lambda)$)
3. Proof (Combination of results proved by
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Minimal affinization

\mathfrak{g} : simple Lie algebra of rank n ,

$\mathcal{L}\mathfrak{g} = \mathfrak{g} \otimes \mathbb{C}[t, t^{-1}]$: loop algebra, $([x \otimes f, y \otimes g] = [x, y] \otimes fg)$

$U_q(\mathcal{L}\mathfrak{g})$: quantum loop algebra/ $\mathbb{C}(q)$ (q -analog of $U(\mathcal{L}\mathfrak{g})$)

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 $U_q(\mathfrak{g})$: quantum group assoc. with \mathfrak{g} (q -analog of $U(\mathfrak{g})$)

Fact (f.d. $U_q(\mathfrak{g})$ -modules)

$$(1) \quad \left\{ \begin{array}{c} \text{f.d. simple } \mathfrak{g}\text{-mod.} \\ \cup \\ V(\lambda) \end{array} \right\} \xleftrightarrow{1:1} P^+ \xleftrightarrow{1:1} \left\{ \begin{array}{c} \text{f.d. simple } U_q(\mathfrak{g})\text{-mod} \\ \cup \\ V_q(\lambda) \end{array} \right\}$$

(2) The cat. of f.d. \mathfrak{g} -modules and $U_q(\mathfrak{g})$ -modules are semisimple.

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Minimal affinization

Fact. V : an arbitrary f.d. simple $U_q(\mathcal{L}\mathfrak{g})$ -module

$\rightsquigarrow \exists! \lambda \in P^+$ s.t. $V \cong V_q(\lambda) \oplus \bigoplus_{\mu < \lambda} V_q(\mu)^{\oplus m_\mu(V)}$ as a $U_q(\mathfrak{g})$ -module.

In this case, V is called an **affinization** of $V_q(\lambda)$.

$\{U_q(\mathfrak{g})\text{-isom. classes of affiniz. of } V_q(\lambda)\} \Leftarrow$ partial order is defined

$([V] \geq [W] \Leftrightarrow \{m_\mu(V)\}_\mu \geq \{m_\mu(W)\}_\mu \text{ w.r.t. lexicographic order})$

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Examples of Minimal affinizations

Minimal affinizations for $\mathfrak{g} = \mathfrak{sl}_{n+1}$

When $\mathfrak{g} = \mathfrak{sl}_{n+1}$, \exists alg. hom. $\varphi: U_q(\mathcal{L}\mathfrak{g}) \rightarrow U_q(\mathfrak{g})$ (evaluation map)
(q -analog of the map $\mathcal{L}\mathfrak{g} \rightarrow \mathfrak{g}: x \otimes f \rightarrow f(a)x$ for any $a \in \mathbb{C}^\times$)

$\rightsquigarrow \varphi^* V_q(\lambda)$: simple $U_q(\mathcal{L}\mathfrak{g})$ -mod. \Leftarrow minimal affinization of $V_q(\lambda)$
($\because \varphi^* V_q(\lambda) \cong V_q(\lambda)$ as a $U_q(\mathfrak{g})$ -mod.)

Remark. If $\mathfrak{g} \neq \mathfrak{sl}_{n+1}$, evaluation map **does not** exist.

\rightsquigarrow Most of minimal affinizations are reducible as a $U_q(\mathfrak{g})$ -module,
and it is not easy to determine the decompositions or characters.

Another example

Kirillov-Reshetikhin modules = minimal affinizations of $V_q(m\overline{\omega}_i)$

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Main Theorem

In the sequel, assume that \mathfrak{g} is of type $ABCD$.

Let $\lambda \in P^+$, and let $L_q(\lambda)$ be a minimal affinization of $V_q(\lambda)$.

Theorem

Assume that
$$\begin{cases} \langle h_n, \lambda \rangle = 0 & \text{if } \mathfrak{g}: \text{ type } BC, \\ \langle h_{n-1}, \lambda \rangle = \langle h_n, \lambda \rangle = 0 & \text{if } \mathfrak{g}: \text{ type } D, \end{cases}$$

and set $\lambda_i := \sum_{k \geq i} \langle h_k, \lambda \rangle \in \mathbb{Z}_{\geq 0}$ for $1 \leq i \leq n$. Then we have

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$$= \begin{cases} \det \left(\text{ch } V((\lambda_i - i + j)\varpi_1) \right)_{1 \leq i, j \leq n} & \mathfrak{g}: ABD \\ \det \left(\sum_{0 \leq 2\ell \leq \lambda_i - i + j} \text{ch } V((\lambda_i - i + j - 2\ell)\varpi_1) \right)_{1 \leq i, j \leq n} & \mathfrak{g}: C \end{cases}$$

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$$= \begin{cases} \det \left(\text{ch } V((\lambda_i - i + j)\varpi_1) \right)_{1 \leq i, j \leq n} & \mathfrak{g}: ABD \\ \det \left(\sum_{0 \leq 2\ell \leq \lambda_i - i + j} \text{ch } V((\lambda_i - i + j - 2\ell)\varpi_1) \right)_{1 \leq i, j \leq n} & \mathfrak{g}: C \end{cases}$$

Remark. In type A , this is JT formula since $\text{ch } L_q(\lambda) = \text{ch } V(\lambda)$.

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$\lambda \in P^+$: as above. For every $\mu \in P^+$,

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Sketch of the proof

Graded limits

$L_q(\lambda): U_q(\mathcal{L}\mathfrak{g})\text{-mod.}/\mathbb{C}(q) \xrightarrow{q \rightarrow 1} L_1(\lambda): \mathcal{L}\mathfrak{g}\text{-mod.}/\mathbb{C}$ (classical limit)

$\xrightarrow{\text{restrict}} L_1(\lambda): \mathfrak{g}[t]\text{-module} \quad (\mathfrak{g}[t] = \mathfrak{g} \otimes \mathbb{C}[t] \subseteq \mathcal{L}\mathfrak{g} = \mathfrak{g} \otimes \mathbb{C}[t, t^{-1}])$

Fact. $\exists a \in \mathbb{C}^\times$ s.t. $(\mathfrak{g} \otimes (t+a)^N) L_1(\lambda) = 0 \quad (N \gg 0)$

\rightsquigarrow Define an auto. τ_a on $\mathfrak{g}[t]$ by $\tau_a(\mathfrak{g} \otimes f(t)) = \mathfrak{g} \otimes f(t+a)$

$L(\lambda) := \tau_a^*(L_1(\lambda))$: **graded limit** of $L_q(\lambda)$ (\mathbb{Z} -graded $\mathfrak{g}[t]$ -module)

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$\mathfrak{g} = \mathfrak{n}_+ \oplus \mathfrak{h} \oplus \mathfrak{n}_-$: triangular decomposition,

Define $\Delta'_+ := \{\alpha \in \Delta_+ \mid \alpha = \sum m_i \alpha_i, m_i \leq 1\} \subseteq \Delta_+$.

Proposition (N)

Let $M(\lambda)$ be the $\mathfrak{g}[t]$ -module generated by a vector v with relations

$$\begin{aligned} \mathfrak{n}_+[t]v &= 0, & (h \otimes t^n)v &= \delta_{0,n} \lambda(h)v \text{ for } h \in \mathfrak{h}, & f_i^{\lambda(h_i)+1}v &= 0, \\ (f_\alpha \otimes t)v &= 0 \text{ for } \alpha \in \Delta'_+. \end{aligned}$$

Then the graded limit $L(\lambda)$ is isomorphic to $M(\lambda)$.

Proposition (Chari-Greenstein, 11)

$$\sum_{(\lambda, s) \in \Gamma(\mu)} (-1)^s \dim \operatorname{Hom}_{\mathfrak{g}}(V(\lambda), \bigwedge^s \mathfrak{g} \otimes V(\mu)) \operatorname{ch} M(\lambda) = \operatorname{ch} V(\mu),$$

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Comment on exceptional types

It would be possible to study minimal affinizations in exceptional types using their graded limits. Indeed, recently we obtain the following polyhedral multiplicity formula for minimal affinizations of type G_2 :

$$L_q(k\varpi_1 + l\varpi_2) \cong_{U_q(\mathfrak{g})} \bigoplus_{(a_1, \dots, a_5) \in S(k, l)} V_q((k - a_1 + a_3 + a_4 - a_5)\varpi_1 + (l - a_2 - 3a_3 - 3a_4)\varpi_2)$$

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